

Non-Private Uniformity and Identity Testing

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Abstract

Re-deriving well-known results in Uniformity Testing

1 Problem Statement

These notes are derived from the excellent survey by Clement Canonne [Can22]. We have n i.i.d samples from an unknown distribution $p \in \text{Simplex}(k)$. We want to guess whether $p = U_k$ or $d_{TV}(p, U_k) \geq \epsilon$, where U_k is the uniform distribution over $\text{Simplex}(k)$. Furthermore, we would like to guess correctly at least two-thirds of the time i.e.

$$\mathbb{P} \left[\text{Guess } p = U_k | p = U_k \right] \geq \frac{2}{3} \quad (1)$$

$$\mathbb{P} \left[\text{Guess } d_{TV}(p, U_k) \geq \epsilon | d_{TV}(p, U_k) \geq \epsilon \right] \geq \frac{2}{3} \quad (2)$$

To guess we need to design some kind of test, that takes as input the parameters k, ϵ and the i.i.d samples and spits out a guess. One very natural test is to count how many of the samples are the same. If the distribution was truly uniformly random, we do not expect to see a lot of overlap between the samples. Before describing the test in the next section we state some general facts about norms and distances.

$$d_{TV}(p, q) = \frac{1}{2} \|p - q\|_1 \leq \frac{\sqrt{k}}{2} \|p - q\|_2 \quad (3)$$

The proof for (3) comes from Cauchy-Schwartz and can be found in any linear algebra textbook under the title “equivalence of norms”. Thus if we have $d_{TV}(p, q) \geq \epsilon$, then we have $\|p - q\|_2 \geq \frac{4\epsilon^2}{k}$. Furthermore, if $q = U_k$, then we have $\|p - U_k\|_2^2 = \|p\|_2^2 - \frac{1}{k}$. Thus combining all these facts we have

- If $d_{TV}(p, U_k) = 0$, then $\|p\|_2^2 = \frac{1}{k}$
- Otherwise, $d_{TV}(p, U_k) \geq \epsilon$, then $\|p\|_2^2 \geq \frac{1+4\epsilon^2}{k}$

Note that $\|p\|_k^k$ is also called the collision probability of k samples drawn from p having the same value.

2 Tester Based on Collisions

Let X_1, \dots, X_n represent n i.i.d samples from p . Let Z denote the number of pairwise collisions i.e

$$Z = \frac{1}{\binom{n}{2}} \sum_{1 \leq s < t \leq n} \mathbb{1}(X_s = X_t) \quad (4)$$

Z represents the fraction of pairwise collisions as a sum of i.i.d indicators and $\mathbb{P}[\mathbb{1}(X_s = X_t)] = \mathbb{E}[\mathbb{1}(X_s = X_t)] = \|p\|_2^2$. Therefore $\mathbb{E}[Z] = \|p\|_2^2$. The obvious test is to pick some threshold τ such that if $Z \leq \tau$, declare that $p = U_k$ and $d_{TV}(p, U_k) \geq \epsilon$ otherwise. Say we pick a threshold between $\frac{1}{k}$

and $\frac{1+4\epsilon^2}{k}$, so say $\tau = \frac{1+2\epsilon^2}{k}$. Then we want to know how many samples we need for equations (1) and (2) to hold.

We need to bound Type I and Type II errors to upper-bound the probability of making a mistake. To upper bound Type I error (or completeness error) we have that the data comes from the uniform distribution but our test threshold still crossed τ , so we need to upper bound the following.

$$\mathbb{P}\left[Z \geq \tau | p = U_k\right] = \mathbb{P}\left[Z \geq \frac{1+2\epsilon^2}{k}\right] \quad (5)$$

$$= \mathbb{P}\left[Z \geq (1+2\epsilon^2)\mathbb{E}[Z]\right] \quad (6)$$

$$\leq \mathbb{P}\left[Z \geq (1+\epsilon^2)\mathbb{E}[Z]\right] \quad (7)$$

We drop the conditional from (5) onwards to make the notation clearer. (6) comes from the assumption that $p = U_k$. To upper bound Type II error, we need upper bound the following quantity

$$\mathbb{P}\left[Z \leq \tau | d_{TV}(p, U_k) \geq \epsilon\right] = \mathbb{P}\left[Z \leq \frac{1+2\epsilon^2}{k}\right] \quad (8)$$

$$\leq \mathbb{P}\left[Z \leq \frac{(1-\epsilon^2)(1+4\epsilon^2)}{k}\right] \quad (9)$$

$$\leq \mathbb{P}\left[Z \leq (1-\epsilon^2)\mathbb{E}[Z]\right] \quad (10)$$

(9) comes from the fact that when $\epsilon \leq \frac{1}{2}$, we have $(1-\epsilon^2)(1+4\epsilon^2) \geq (1+2\epsilon^2)$. (10) comes from the fact that $\mathbb{E}[Z] \geq \frac{(1+4\epsilon^2)}{k}$ from our assumption. Thus if we combined equation (10) and (7) and applied Chebychev we would get

$$\mathbb{P}\left[|Z - \mathbb{E}[Z]| \leq \epsilon^2 \mathbb{E}[Z]\right] \leq \frac{\text{Var}(Z)}{\epsilon^4 \mathbb{E}[Z]^2} \quad (11)$$

If the variance of the estimator is much larger than the wiggle room we've given us, then we would have no hope i.e. we would like $\sqrt{\text{Var}(Z)} \ll (\frac{1+4\epsilon^2}{k} - \frac{1}{k})$.

2.1 Simple but non-optimal upper bound for variance

$$\text{Var}(Z) = \mathbb{E}[Z^2] - \mathbb{E}[Z]^2 \quad (12)$$

We already have $\mathbb{E}[Z] = \|p\|_2^2$, so $\mathbb{E}[Z]^2 = \|p\|_2^4$. We need a manageable expression for $\mathbb{E}[Z^2]$.

$$\mathbb{E}[Z^2] = \frac{1}{\binom{n}{2}^2} \mathbb{E}\left[\sum_{1 \leq s < t \leq n} \sum_{1 \leq s' < t' \leq n} \mathbf{1}(X_s = X_t) \mathbf{1}(X_{s'} = X_{t'})\right] \quad (13)$$

$$= \frac{1}{\binom{n}{2}^2} \sum_{1 \leq s < t \leq n} \sum_{1 \leq s' < t' \leq n} \mathbb{E}[\mathbf{1}(X_s = X_t) \mathbf{1}(X_{s'} = X_{t'})] \quad (14)$$

$$(15)$$

(14) : By Linearity of expectation. Now looking at all those indices in that awkward-looking expression $\mathbb{E}[\mathbf{1}(X_s = X_t) \mathbf{1}(X_{s'} = X_{t'})]$, we can group them into 3 categories based on the cardinality of the set $\{s, t, s', t'\}$ (1) When all the indices are distinct (2) When both pairs are the same (3) When there are 3 distinct indices. Additionally, we will only need to consider the events when the product of the two terms is 1.

1. s, t, s', t' are all distinct i.e. $|\{s, t, s', t'\}| = 4$. Therefore the two collisions are independent of each other. Therefore we get $\mathbb{E}[\mathbf{1}(X_s = X_t) \mathbf{1}(X_{s'} = X_{t'})] = \|p\|_2^4$. There are $\binom{n}{2} \binom{n-2}{2} = 6 \binom{n}{4}$ ways to select pairs such they are distinct.
2. $s = s'$ and $t = t'$ i.e. both pairs are the same $|\{s, t, s', t'\}| = 2$, as the square of an indicator function is just the value of the indicator we have $\mathbb{E}[\mathbf{1}(X_s = X_t) \mathbf{1}(X_{s'} = X_{t'})] = \|p\|_2^2$. There are $\binom{n}{2}$ ways of selecting 2 unique indices.

3. Finally the mankiest event, where we have $|\{s, t, s', t'\}| = 3$. Three distinct indices will take up the same value and one of the pairs will have a repeat index. This is the same as saying 3 items will collide and the probability of having 3 collisions is $\|p\|_3^3$. There are $\binom{n}{3}$ ways of picking three distinct indices. Then for each of those three indices, there are 2 choices for which the remaining index is duplicated. Thus the total number of options is $(2 \times 3)\binom{n}{3}$. Also by the monotonicity of norms, we have $\|p\|_3^3 \leq \|p\|_2^3$ which we will use below.

Thus we have

$$\text{Var}(Z) = \left[\frac{1}{\binom{n}{2}^2} \sum_{1 \leq s < t \leq n} \sum_{1 \leq s' < t' \leq n} \mathbb{E}[\mathbb{1}(X_s = X_t) \mathbb{1}(X_{s'} = X_{t'})] \right] - \mathbb{E}[Z]^2 \quad (16)$$

$$= \frac{1}{\binom{n}{2}^2} \left[6 \binom{n}{4} \|p\|_2^4 + \binom{n}{2} \|p\|_2^2 + 6 \binom{n}{3} \|p\|_3^3 \right] - \frac{\binom{n}{2}^2}{\binom{n}{2}^2} \|p\|_2^4 \quad (17)$$

$$= \frac{1}{\binom{n}{2}^2} \left[\|p\|_2^4 \left(6 \binom{n}{4} - \binom{n}{2}^2 \right) + \binom{n}{2} \|p\|_2^2 + 6 \binom{n}{3} \|p\|_3^3 \right] \quad (18)$$

$$\leq \frac{1}{\binom{n}{2}^2} \left[\|p\|_2^4 \left(6 \binom{n}{4} - \binom{n}{2}^2 \right) + \binom{n}{2} \|p\|_2^2 + 6 \binom{n}{3} \|p\|_2^3 \right] \quad (19)$$

$$\leq \frac{1}{\binom{n}{2}^2} \left[\binom{n}{2} \|p\|_2^2 + 6 \binom{n}{3} \|p\|_2^3 \right] \quad (20)$$

$$= \frac{1}{\binom{n}{2}^2} \left[\binom{n}{2} \mathbb{E}[Z] + 6 \binom{n}{3} \mathbb{E}[Z]^{3/2} \right] \quad (21)$$

$$\leq O\left(\frac{1}{n^2}\right) \mathbb{E}[Z] + O\left(\frac{1}{n}\right) \mathbb{E}[Z]^{3/2} \quad (22)$$

(19): Monotonicity, $\|p\|_3^3 \leq \|p\|_2^3$

(20) $\left(6 \binom{n}{4} - \binom{n}{2}^2 \right) \leq 0$, for $n \geq 2$

(22): $\frac{4}{n^2} \geq \frac{1}{\binom{n}{2}}$ and

Plugging this into (11) we have

$$\mathbb{P} \left[|Z - \mathbb{E}[Z]| \leq \epsilon^2 \mathbb{E}[Z] \right] \leq \frac{O\left(\frac{1}{n^2}\right) \mathbb{E}[Z] + O\left(\frac{1}{n}\right) \mathbb{E}[Z]^{3/2}}{\epsilon^4 \mathbb{E}[Z]^2} \quad (23)$$

$$= \frac{1}{\epsilon^4 n^2 \mathbb{E}[Z]} + \frac{1}{\epsilon^4 n \mathbb{E}[Z]^{1/2}} \quad (24)$$

$$\leq \frac{k}{\epsilon^4 n^2} + \frac{\sqrt{k}}{\epsilon^4 n} \quad (25)$$

(25): $\mathbb{E}[Z] \geq \frac{1}{k}$ for both cases.

Setting (25) equal to 1/3 and solving for n we get

$$n \geq O\left(\frac{\sqrt{k}}{\epsilon^4}\right) \quad (26)$$

2.2 Optimal Upper Bound

Unfortunately, this bound can be improved to

$$n \geq O\left(\frac{\sqrt{k}}{\epsilon^2}\right) \quad (27)$$

Thus we must account for some slack in the analysis for variance. Thus going back to our analysis where we kicked out the negative term

$$\binom{n}{2}^2 \text{Var}(Z) = \left[\|p\|_2^4 \left(6 \binom{n}{4} - \binom{n}{2}^2 \right) + \binom{n}{2} \|p\|_2^2 + 6 \binom{n}{3} \|p\|_3^3 \right] \quad (28)$$

$$= \left[\|p\|_2^4 \left(-6 \binom{n}{3} - \binom{n}{2} \right) + \binom{n}{2} \|p\|_2^2 + 6 \binom{n}{3} \|p\|_3^3 \right] \quad (29)$$

$$= \binom{n}{2} \|p\|_2^2 (1 - \|p\|_2^2) + 6 \binom{n}{3} (\|p\|_3^3 - \|p\|_4^2) \quad (30)$$

$$(29) : \binom{n}{2}^2 = 6 \binom{n}{4} + 6 \binom{n}{3} + \binom{n}{2}$$

The first summand is the variance of a binomial distribution with parameters $\text{Bin}(\binom{n}{2}, \|p\|_2^2)$ and the second summand is always non-negative (where the middle inequality comes from cauchy-schwartz)

$$\|p\|_2^4 = \left(\sum_i p(i)^{3/2} p(i)^{1/2} \right)^2 \leq \sum_i p(i)^3 \sum_i p(i) = \sum_i p(i)^3 = \|p\|_3^3 \quad (31)$$

Thus going back to our variance analysis and dividing both sides by $\binom{n}{2}^2$ we get

$$\text{Var}(Z) = \frac{1}{\binom{n}{2}^2} \left[\binom{n}{2} \|p\|_2^2 (1 - \|p\|_2^2) + 6 \binom{n}{3} (\|p\|_3^3 - \|p\|_4^2) \right] \quad (32)$$

$$\leq \frac{1}{\binom{n}{2}^2} \left[\binom{n}{2} \|p\|_2^2 + 6 \binom{n}{3} (\|p\|_3^3 - \|p\|_4^2) \right] \quad (33)$$

$$\leq O\left(\frac{1}{n^2}\right) \|p\|_2^2 + O\left(\frac{1}{n}\right) (\|p\|_3^3 - \|p\|_4^2) \quad (34)$$

Now if $p = U_k$, then $\text{Var}(Z) = O\left(\frac{1}{n^2}\right) \mathbb{E}[Z]$ since $\|p\|_3^3 = \|p\|_4^2$, thus plugging it back to (11), we get

$$\mathbb{P}[Z \geq \tau] \leq \mathbb{P}\left[Z \geq (1 + \epsilon^2) \mathbb{E}[Z]\right] \leq \frac{k}{n^2 \epsilon^4} \quad (35)$$

Setting the RHS of (35) equal to 1/3 we get

$$n \geq O\left(\frac{\sqrt{k}}{\epsilon^2}\right) \quad (36)$$

Okay but this only solves the case for Type I error, we still need to deal with the case that $d_{TV}(p, U_k) \geq \epsilon$. Define $\alpha^2 := k \|p - U_k\|_2^2 \geq 4\epsilon^2$, then $\mathbb{E}[Z_1] = \frac{1+\alpha^2}{k}$.

$$\mathbb{P}[Z \leq \tau] = \mathbb{P}\left[Z \leq \frac{1+2\epsilon^2}{k}\right] \quad (37)$$

$$= \mathbb{P}\left[Z \leq \frac{1+2\epsilon^2}{1+\alpha^2} \mathbb{E}[Z]\right] \quad (38)$$

$$= \mathbb{P}\left[Z \leq \left(1 - \frac{\alpha^2 - 2\epsilon^2}{1+\alpha^2}\right) \mathbb{E}[Z]\right] \quad (39)$$

$$\leq \mathbb{P}\left[Z \leq \left(1 - \frac{\alpha^2}{2(1+\alpha^2)}\right) \mathbb{E}[Z]\right] \quad (40)$$

$$\leq \frac{\text{Var}(Z)}{\alpha^4 \mathbb{E}[Z]^2} 4(1+\alpha^2)^2 \quad (41)$$

$$\leq \frac{16(1+\alpha^2)^2 \mathbb{E}[Z]}{\alpha^4 \mathbb{E}[Z]^2} \frac{1}{n^2} + \frac{16(1+\alpha^2)^2}{\alpha^4 n \mathbb{E}[Z]^2} (\|p\|_3^3 - \|p\|_4^2) \quad (42)$$

(38): How we defined α , (39): By the assumption that p is far from U_k , (40): $\alpha^2 \geq 4\epsilon^2$
Dealing with the left summand first. We have $\mathbb{E}[Z] = \frac{1+\alpha^2}{k}$. Thus we get

$$\frac{16(1+\alpha^2)^2}{\alpha^4 \mathbb{E}[Z]^2} \frac{\mathbb{E}[Z]}{n^2} = \frac{16(1+\alpha^2)k}{\alpha^4 n^2} \leq \frac{5k}{\epsilon^4 n^2} \quad (43)$$

(43): For $x > 0$, $\frac{1+x}{x^2}$ is decreasing function. Furthermore $\alpha^2 \geq 4\epsilon^2$ and $\epsilon \leq 1$, thus $1+4\epsilon \leq 5$

Ok so the one last thing to show is the last summand $\frac{4}{n\mathbb{E}[Z]^2} \left(\|p\|_3^3 - \|p\|_4^2 \right)$ and then we would be done. As $\|p\|_2^2 \geq \frac{1}{k}$ we have

$$\left(\|p\|_3^3 - \|p\|_4^2 \right) \leq \|p - U_k + U_k\|_3^3 - \frac{1}{k^2} \quad (44)$$

$$= \|p - U_k\|_3^3 + \frac{3}{k} \|p - u\|_2^2 \quad (45)$$

$$\leq \|p - U_k\|_2^3 + \frac{3}{k} \|p - u\|_2^2 \quad (46)$$

$$= \frac{\alpha^3}{k^{3/2}} + \frac{3\alpha^2}{k^2} \quad (47)$$

Finally, multiplying it with $\frac{16(1+\alpha^2)^2}{n\alpha^4 \|p\|_2^4}$

$$\frac{16(1+\alpha^2)^2}{n\alpha^4 \|p\|_2^4} \left(\|p\|_3^3 - \|p\|_4^2 \right) \leq \frac{16(1+\alpha^2)^2}{n\alpha^4 \|p\|_2^4} \frac{\alpha^3}{k^{3/2}} + \frac{3\alpha^2}{k^2} \quad (48)$$

$$\leq \frac{8\sqrt{k}}{\epsilon n} + \frac{12}{n\epsilon^2} \quad (49)$$

Combining the two summands, we get

$$n \geq O\left(\frac{\sqrt{k}}{\epsilon^2}\right) \quad (50)$$

References

- [Can22] Clément L. Canonne. Topics and Techniques in Distribution Testing: A Biased but Representative Sample. March 2022.